

FuSe

an(other) OCaml implementation
of binary sessions

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target API

Linear type theory for asynchronous session types (Gay & Vasconcelos 2010)

accept : $A \text{ service} \rightarrow A$

request : $A \text{ service} \rightarrow \bar{A}$

send : $\alpha \rightarrow !\alpha.A \rightarrow A$

receive : $?\alpha.A \rightarrow \alpha \times A$

select : $(\bar{A}_k \rightarrow [C_i : \bar{A}_i]_{i \in I}) \rightarrow \oplus [C_i : A_i]_{i \in I} \rightarrow A_k$

branch : $\& [C_i : A_i]_{i \in I} \rightarrow [C_i : A_i]_{i \in I}$

create : $\text{unit} \rightarrow A \times \bar{A}$

close : $\text{end} \rightarrow \text{unit}$

duality and linearity in OCaml

Turn duality into equality

- ▶ Dardha et al. 2012 + tweak

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Ostrich approach to linearity

- ▶ no type-based mechanism, no monads, ...
- ▶ affinity violations detected at runtime (Tov & Pucella 2010, Hu & Yoshida 2016)
- ▶ some (many?) violations detected by OCaml anyway



duality \Rightarrow equality

Session types revisited

(Dardha et al., PPDP 2012)

$$\llbracket T \rrbracket = \kappa[t]$$

$$\llbracket !t.T \rrbracket = ![t \times \llbracket \bar{T} \rrbracket]$$

$$\llbracket ?t.T \rrbracket = ?[t \times \llbracket T \rrbracket]$$

$$\llbracket T \rrbracket = \kappa[t] \Rightarrow \llbracket \bar{T} \rrbracket = \bar{\kappa}[t]$$

duality \Rightarrow equality

Session types revisited + tweak

(Dardha et al., PPDP 2012)

$$\llbracket T \rrbracket = \langle t, s \rangle$$

$$\llbracket !t.T \rrbracket = \langle 0, t \times \llbracket \bar{T} \rrbracket \rangle$$

$$\llbracket ?t.T \rrbracket = \langle t \times \llbracket T \rrbracket, 0 \rangle$$

$$\llbracket T \rrbracket = \langle t, s \rangle \Rightarrow \llbracket \bar{T} \rrbracket = \langle s, t \rangle$$

target API (encoded)

accept : $\langle \alpha, \beta \rangle \text{ service} \rightarrow \langle \alpha, \beta \rangle$
request : $\langle \alpha, \beta \rangle \text{ service} \rightarrow \langle \beta, \alpha \rangle$

send : $\gamma \rightarrow \langle \mathbb{0}, \gamma \times \langle \beta, \alpha \rangle \rangle \rightarrow \langle \alpha, \beta \rangle$
receive : $\langle \gamma \times \langle \alpha, \beta \rangle, \mathbb{0} \rangle \rightarrow \gamma \times \langle \alpha, \beta \rangle$

select : $(\langle \beta, \alpha \rangle \rightarrow \gamma) \rightarrow \langle \mathbb{0}, \gamma \rangle \rightarrow \langle \alpha, \beta \rangle$
branch : $\langle \gamma, \mathbb{0} \rangle \rightarrow \gamma$

create : **unit** $\rightarrow \langle \alpha, \beta \rangle \times \langle \beta, \alpha \rangle$
close : $\langle \mathbb{0}, \mathbb{0} \rangle \rightarrow \text{unit}$

DEMO

wrap-up

- ▶ GV-style sessions (Gay & Vasconcelos, JFP 2010)
- ▶ context-free session types (Thiemann & Vasconcelos, ICFP 2016)
- ▶ subtyping (Gay & Hole, Acta 2005)
- ▶ chaperone contracts (ongoing with Hernán Melgratti)

References

- ▶ FuSe details and implementation (Padovani, JFP 2017)
- ▶ context-free session types in FuSe (Padovani, ESOP 2017)