

Inference of Global Progress Properties for Dynamically Interleaved Multiparty Sessions

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BEAT 2013

The problem

$$a(y).b(z).y?(x).z! \langle x \rangle$$
$$\bar{a}(y).\bar{b}(z).z?(x).y! \langle x \rangle$$

- two distinct sessions
- each session is well typed
- the system gets stuck

The problem

$$a(y).b(z).y?(x).z!(x) \quad y : ?\text{int} \quad z : !\text{int}$$
$$\bar{a}(y).\bar{b}(z).z?(x).y!(x) \quad y : !\text{int} \quad z : ?\text{int}$$

- two distinct sessions
- each session is well typed
- the system gets stuck

The “interaction” type system

If $\vdash P$, then P never gets stuck

- 😊 Bettini, Coppo, D'Antoni, De Luca, Dezani-Ciancaglini, Yoshida,
Global Progress in Dynamically Interleaved Multiparty Sessions, CONCUR 2008
- 😢 **not syntax-directed**

Outline

- ① Progress
- ② Key ideas of the type system
- ③ Two examples
- ④ Remarks

Progress 1/2

If $P \rightarrow^* \mathcal{E} [s?(x).P']$
Then $\rightarrow^* \mathcal{E}'[s?(x).P' \mid s : m \cdot h]$

If $P \rightarrow^* \mathcal{E} [s : m \cdot h]$
Then $\rightarrow^* \mathcal{E}'[s : m \cdot h \mid s?(x).P']$

A process without progress

$$a(y).b(z).y?(x).z! \langle x \rangle \mid \bar{a}(y).\bar{b}(z).z?(x).y! \langle x \rangle$$

A process without progress

$$a(y).b(z).y?(x).z! \langle x \rangle \mid \bar{a}(y).\bar{b}(z).z?(x).y! \langle x \rangle$$

\downarrow_*

$$(\nu s)(\nu s')(s?(x).s'! \langle x \rangle \mid s'? (x).s! \langle c \rangle \mid s : \emptyset \mid s' : \emptyset)$$

A process without progress

$$a(y).b(z).y?(x).z! \langle x \rangle \mid \bar{a}(y).\bar{b}(z).z?(x).y! \langle x \rangle$$
 \downarrow_{*}
$$(\nu s)(\nu s') (s?(x).s'! \langle x \rangle \mid s'? \langle x \rangle .s! \langle c \rangle \mid s : \emptyset \mid s' : \emptyset)$$

Progress 2/2

A **good** process that looks like a **bad** one

$$P \rightarrow^* \mathcal{E}[s?(x).P' \mid \overline{b}(y).s!(3).Q']$$

A **bad** process that looks like a **good** one

$c(y)$.(a process that gets stuck)

Progress 2/2

A **good** process that looks like a **bad** one

$$P \rightarrow^* \mathcal{E}[s?(x).P' \mid \overline{b}(y).s!(3).Q']$$

A **bad** process that looks like a **good** one

$c(y).$ (a process that gets stuck)

Idea

- define progress modulo **catalyzers**
- catalyst = missing participant that never gets stuck

Consequence

- session initiation can be considered **non-blocking**

Interaction type system: basic ideas

- ① associate processes with **dependencies** $a \prec b$

“an action of service a blocks an action of service b ”

- ② a process is well typed if it yields **no circular dependencies**

Computing service dependencies

$$a(y).b(z).y?(x).z! \langle x \rangle \quad a \prec b$$

$$\bar{a}(y).\bar{b}(z).z?(x).y! \langle x \rangle \quad b \prec a$$

Service names as messages

$$a(y).b(z).y?(x).z! \langle x \rangle \quad a \prec b$$

$$\bar{c}(t).t?(x).\bar{x}(y).\bar{b}(z).z?(x).y! \langle x \rangle$$

$$c(t).t! \langle a \rangle$$

Service names as messages

$$a(y).b(z).y?(x).z!(x) \quad a \prec b$$

$$t?(x).\bar{x}(y).\bar{b}(z).z?(x).y!(x)$$

$$t!(a)$$



Service names as messages

$$a(y).b(z).y?(x).z!(x) \quad a \prec b$$

$$\bar{a}(y).\bar{b}(z).z?(x).y!(x) \quad b \prec a$$

Service names as messages

$$a(y).b(z).y?(x).z!(x) \qquad a \prec b$$

$$\bar{a}(y).\bar{b}(z).z?(x).y!(x)$$

Idea

- identify a class of safe services even if mutually dependent
- restrict messages to services in this class

Nested services

Definition

a is a **nested service** if $\lambda \prec a$ implies that λ is a nested service

Nested?

$$\bar{a}(y).\bar{a}(z).z?(x).y?(x')$$
 $a \prec a$ ✓

$$\begin{array}{l} \bar{a}(y).\bar{b}(z).z?(x).y?(x') \\ | \quad \bar{b}(z).\bar{a}(y).y?(x).z?(x') \end{array}$$
 $b \prec a$ ✓
 $a \prec b$

$$\bar{a}(y).\bar{b}(z).y?(x).z?(x')$$
 $y \prec b$ ✗

Private services

$$a(y).(\nu b)(b(z).z?(x).y!(\langle x \rangle))$$

- no catalyst can help starting the session on b

Private services

$$a(y).(\nu b)(b(z).z?(x).y!(x))$$

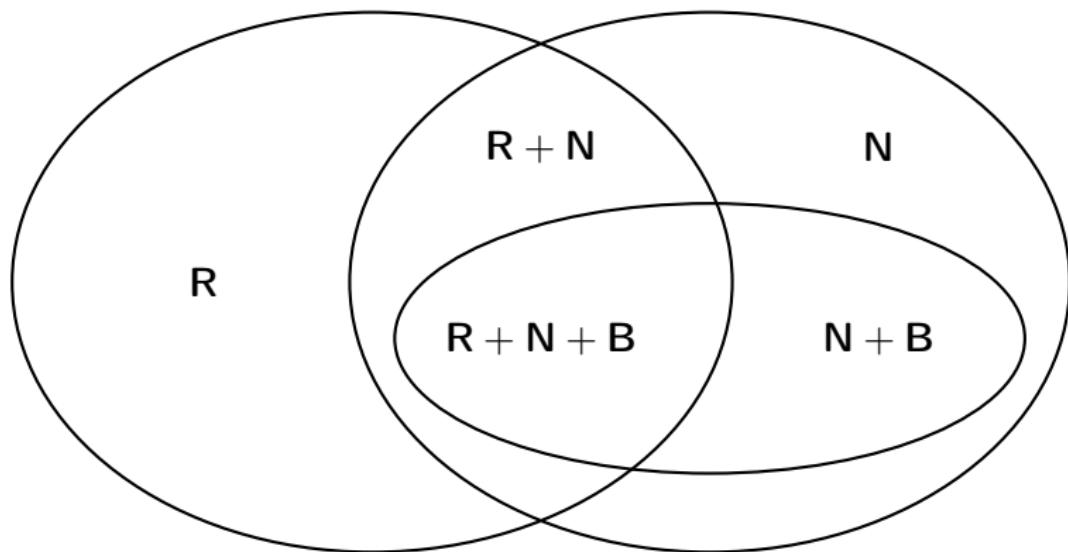
- no catalyst can help starting the session on b

Definition

a is **boundable** if it is never followed by free channels

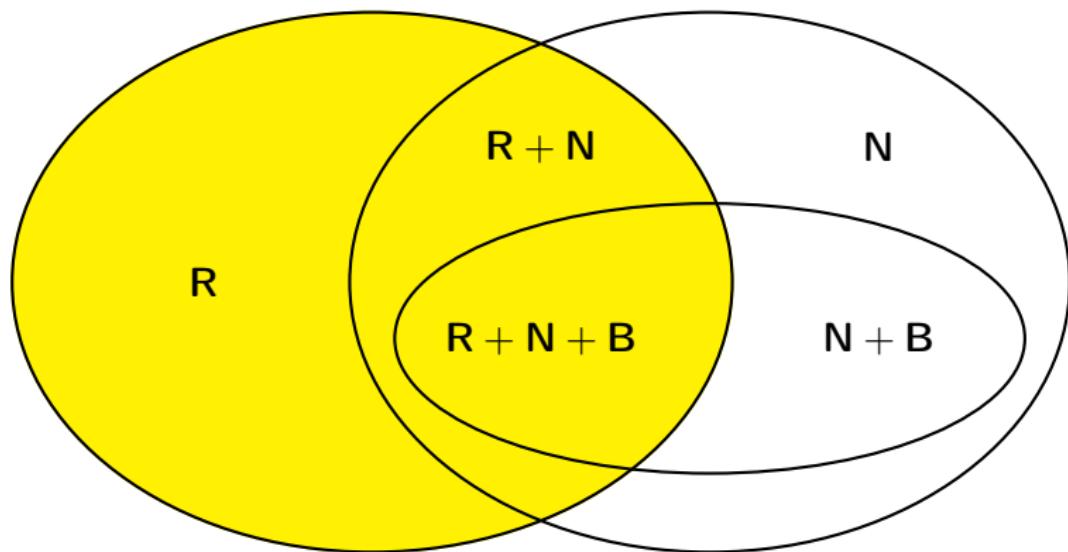
- b is nested but not boundable

Service classification



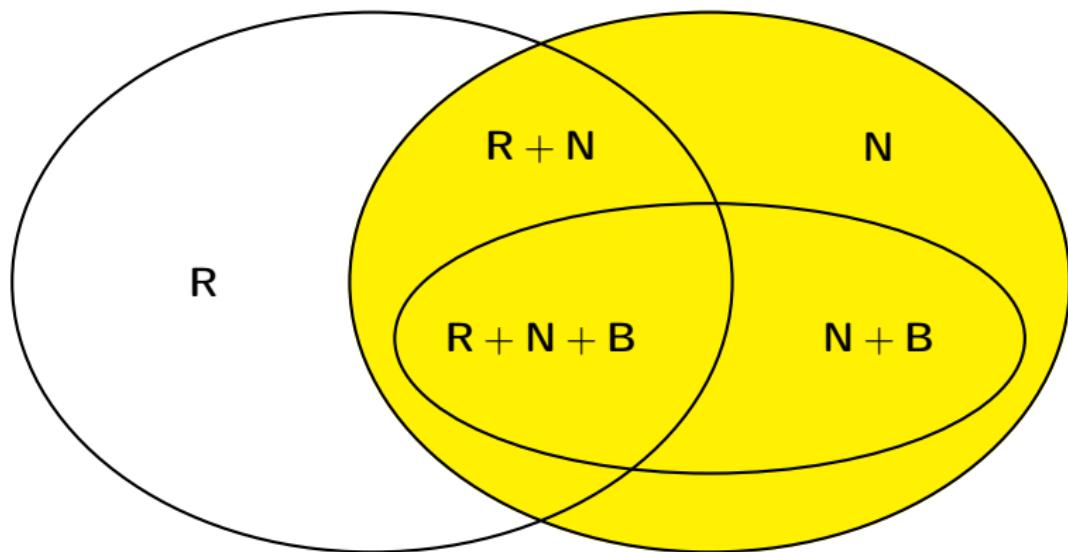
- each service can have up to three features ...
- ... which the interaction type system **guesses**

Service classification



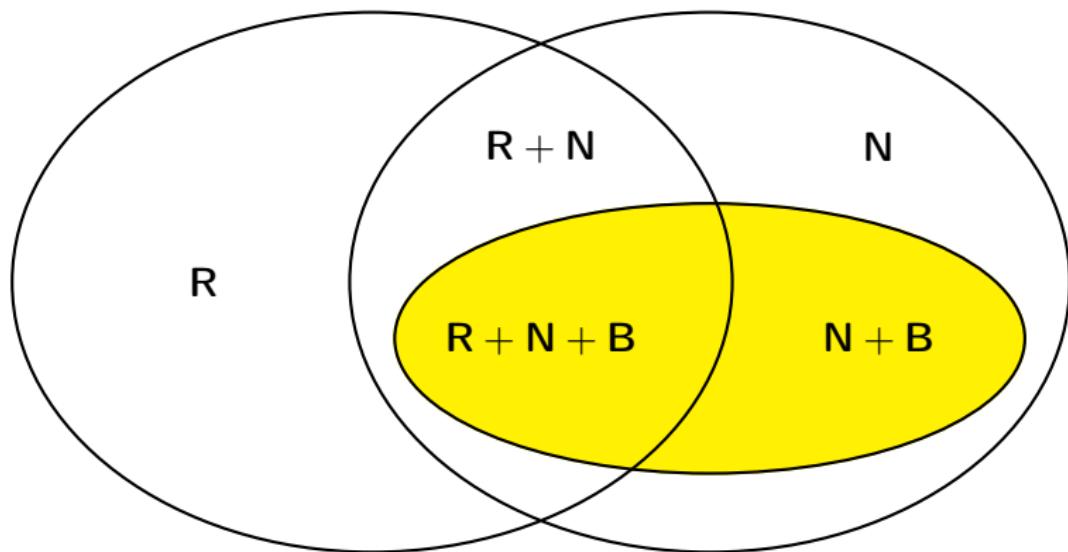
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Service classification



- each service can have up to three features ...
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Service classification



- each service can have up to three features ...
- ... which the interaction type system **guesses**

Algorithm judgments

$P \Rightarrow D; R; N; B$

If...

$$\begin{aligned}D^\infty &\subseteq N \setminus R \\D \downarrow N &\subseteq N \\fs(P) &\subseteq R \cup N\end{aligned}$$

Example 1

$$a(y).b(z).y?(x).z! \langle x \rangle \Rightarrow$$

Example 1

$$\frac{\frac{0 \Rightarrow}{z! \langle x \rangle \Rightarrow} \frac{}{y? \langle x \rangle . z! \langle x \rangle \Rightarrow} \frac{}{b(z) . y? \langle x \rangle . z! \langle x \rangle \Rightarrow} \frac{}{a(y) . b(z) . y? \langle x \rangle . z! \langle x \rangle \Rightarrow}}$$

Example 1

all services have all features

$$\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{z!(x) \Rightarrow}{\frac{y?(x).z!(x) \Rightarrow}{\frac{b(z).y?(x).z!(x) \Rightarrow}{\frac{a(y).b(z).y?(x).z!(x) \Rightarrow}}}}}$$

Example 1

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{z!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{y?(x).z!(x) \Rightarrow} \\ \frac{}{b(z).y?(x).z!(x) \Rightarrow} \\ \frac{}{a(y).b(z).y?(x).z!(x) \Rightarrow}$$

Example 1

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{z!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\frac{y?(x).z!(x) \Rightarrow \{y \prec z\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{b(z).y?(x).z!(x) \Rightarrow}{a(y).b(z).y?(x).z!(x) \Rightarrow}}}$$

Example 1

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{z!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{y?(x).z!(x) \Rightarrow \{y \prec z\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$
$$\frac{\frac{y?(x).z!(x) \Rightarrow \{y \prec z\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{b(z).y?(x).z!(x) \Rightarrow \{y \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}}{a(y).b(z).y?(x).z!(x) \Rightarrow}$$

$B \subseteq N$

$D \downarrow N \subseteq N$

Example 1

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{z!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\frac{y?(x).z!(x) \Rightarrow \{y \prec z\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{b(z).y?(x).z!(x) \Rightarrow \{y \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}{a(y).b(z).y?(x).z!(x) \Rightarrow \{\textcolor{red}{a} \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}}}$$

Example 1 (cont.)

 $\overline{a}(y).\overline{b}(z).z?(x).y!\langle x \rangle$

 $a(y) \cdots | \overline{a}(y) \cdots \Rightarrow$

Example 1 (cont.)

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{y! \langle x \rangle \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\frac{z? (x). y! \langle x \rangle \Rightarrow \{z \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{\bar{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\bar{a}(y). \bar{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}}$$

$a(y) \cdots | \bar{a}(y) \cdots \Rightarrow$

Example 1 (cont.)

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{y! \langle x \rangle \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\frac{z? (x). y! \langle x \rangle \Rightarrow \{z \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{\bar{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\bar{a}(y). \bar{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}}$$

$$a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}$$

$$a(y) \cdots \mid \bar{a}(y) \cdots \Rightarrow$$

Example 1 (cont.)

$$\frac{\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{y! \langle x \rangle \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{z? (x). y! \langle x \rangle \Rightarrow \{z \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\overline{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$
$$\frac{\overline{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\overline{a}(y). \overline{b}(z). z? (x). y! \langle x \rangle \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$

$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{a}(y) \cdots \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \overline{a}(y) \cdots \Rightarrow}$$

Example 1 (cont.)

$$\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{y!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\frac{z?(x).y!(x) \Rightarrow \{z \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{\bar{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{\bar{a}(y).\bar{b}(z).z?(x).y!(x) \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \text{D}^\infty \subseteq \mathbb{N} \setminus \mathbb{R} \quad \text{a} \prec b \Rightarrow \{a \prec b, b \prec a\}; \mathcal{S} \setminus \{a, b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}} \quad \text{a} \prec b \Rightarrow \{a \prec b, b \prec a\}; \mathcal{S} \setminus \{a, b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}}$$

Example 1 (cont.)

$$\frac{\frac{\frac{0 \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}{y!(x) \Rightarrow \emptyset; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{z?(x).y!(x) \Rightarrow \{z \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$
$$\frac{\overline{b}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{a}(y) \cdots \Rightarrow \{b \prec a\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \overline{a}(y) \cdots \Rightarrow \{a \prec b, b \prec a\}; \mathcal{S} \setminus \{a, b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}}$$

Example 2

$$\overline{\bar{c}(t).t?(x).\bar{x}(y).\bar{b}(z).z?(x).y! \langle x \rangle} \Rightarrow$$

$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \bar{c}(t) \cdots \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \bar{c}(t) \cdots \Rightarrow \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}$$

Example 2

$$\frac{\overline{\vdots} \quad \overline{\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\overline{x(y).\overline{b}(z).z?(x).y!(x) \Rightarrow} \quad \overline{\overline{t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow}} \quad \overline{\overline{\overline{c}(t).t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow}}}$$
$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{c}(t) \cdots \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \overline{c}(t) \cdots \Rightarrow \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}$$

Example 2

$$\frac{\frac{\frac{\vdots}{x \text{ must be nested, so } b \text{ too}}}{\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}}{t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \text{dependencies discharged}} \overline{c}(t).t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow$$

$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{c}(t) \cdots \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}{a(y) \cdots | \overline{c}(t) \cdots \Rightarrow \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}$$

Example 2

$$\frac{\vdots}{\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$
$$\frac{}{\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}$$
$$\frac{}{t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}$$
$$\frac{}{\overline{c}(t).t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}$$
$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{c}(t) \cdots \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \overline{c}(t) \cdots \Rightarrow \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}$$

Example 2

$$\frac{\vdots}{\overline{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}$$
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$$\frac{}{\overline{c}(t).t?(x).\overline{x}(y).\overline{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}$$

$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{c}(t) \cdots \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}{a(y) \cdots \mid \overline{c}(t) \cdots \Rightarrow \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}$$

Example 2

$$\frac{\overline{\bar{b}(z).z?(x).y!(x) \Rightarrow \{b \prec y\}; \mathcal{S}; \mathcal{S}; \mathcal{S}}}{\overline{\bar{x}(y).\bar{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}}$$
$$\frac{\overline{t?(x).\bar{x}(y).\bar{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}}{\overline{\bar{c}(t).t?(x).\bar{x}(y).\bar{b}(z).z?(x).y!(x) \Rightarrow \emptyset; \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}}$$

$$\frac{a(y) \cdots \Rightarrow \{a \prec b\}; \mathcal{S}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\} \quad \overline{\bar{c}(t) \cdots \mapsto \emptyset, \mathcal{S} \setminus \{b\}; \mathcal{S}; \mathcal{S}}}{a(y) \cdots \mid \overline{\bar{c}(t) \cdots \mapsto \{a \prec b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}; \mathcal{S} \setminus \{b\}}}$$

Result

Theorem

If $P \Rightarrow D; R; N; B$, then P has progress

Proof.

The algorithm is sound and complete wrt the inference type system
(cf. CONCUR 2008)



Result

Theorem

If $P \Rightarrow D; R; N; B$, then P has progress

Proof.

The algorithm is sound and complete wrt the inference type system
(cf. CONCUR 2008) (for finite processes only)



Soon to come

Inference for recursive processes

Wrap up

- static analysis for (multiparty) session interleaving
- progress \neq absence of deadlock
 - diverging systems do not necessarily have progress
 - catalyzers may help reduction
- efficient inference algorithm

Future work

- many simple program patterns are **ill typed**
 - 👉 more flexible type discipline is required
- π -calculus \neq **programming language**
 - 👉 richer/more compositional types are needed
- type systems for liveness properties are **complex**
 - 👉 traditional concepts/techniques (fairness, subtyping, coinductive reasoning, ...) must be revisited

A simple ill-typed process with progress

```
def X(y,z) = y!⟨3⟩.z?(x).X(y,z) in      b ⊢ a
def Y(y,z) = y?(x).z!⟨x⟩.Y(y,z) in      a ⊢ b
a(y).b(z).X(y,z) | ¯a(y). ¯b(z).Y(y,z)
```